Virtual Memory

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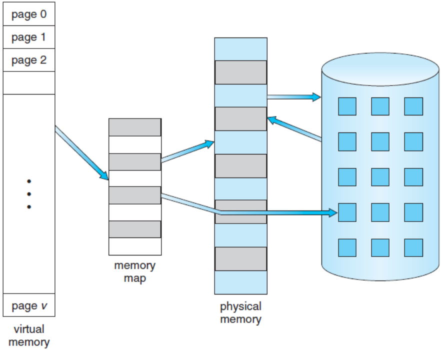
- Only part of the program needs to be in memory for execution.
- More programs running concurrently.

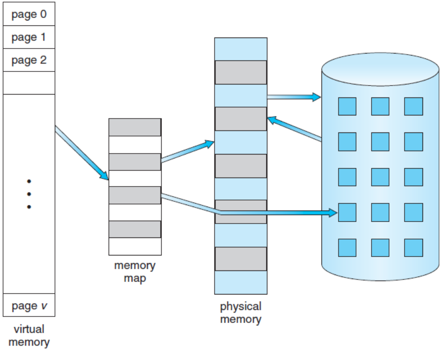
Virtual Memory

- Separation of user logical memory from physical memory.



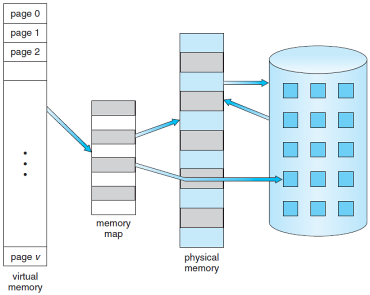
Virtual Memory

- Separation of user logical memory from physical memory.
- Logical address space can be much larger than physical address space.



Virtual Address Space (1/2)

- Virtual address space: logical view of how process is stored in memory.



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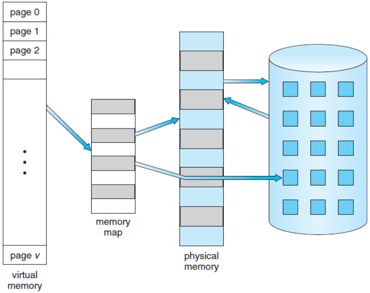

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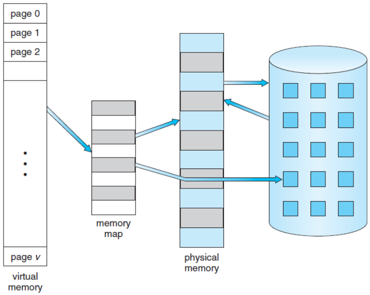
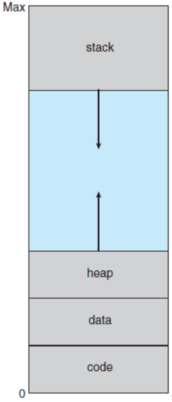
- MMU must map logical to physical.



Virtual Address Space (2/2)

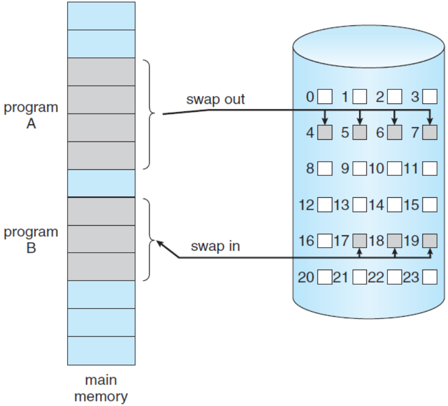
- The hole between heap and stack is part of the virtual address space, but will require actual physical pages only if the heap or stack grows.



Demand Paging

Demand-Paging

- Demand-paging: bring a page into memory only when it is needed.
- lazy swapper.



Basic Concepts

- The pager guesses which pages will be used again before swapping out.

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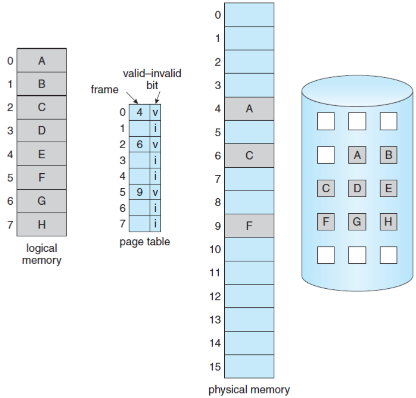
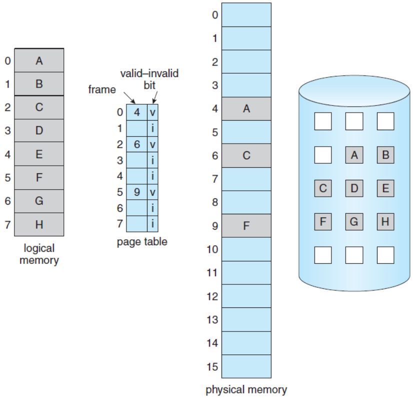


Basic Concepts

- The pager guesses which pages will be used again before swapping out.
- Valid-invalid bit: distinguish between the pages in memory and on disk.
  v: memory resident
  i: not in memory

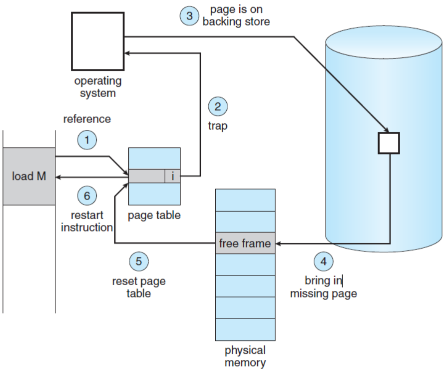


Page Fault

- Access to a page marked invalid causes a page fault.
- Causing a trap to the OS: brings the desired page into memory.

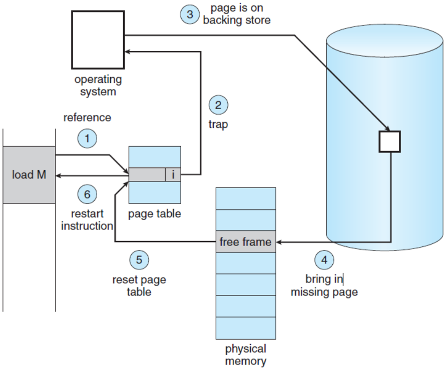
Handling Page Fault (1/6)

- Check an internal table for the process to determine whether the reference was a valid or an invalid memory access.



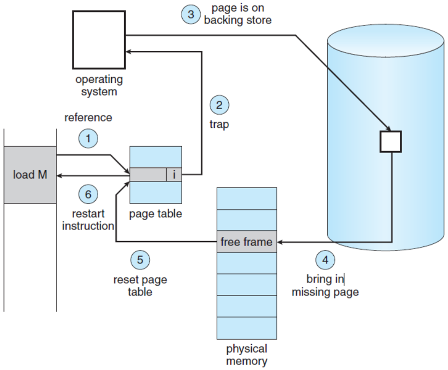
Handling Page Fault (2/6)

- If the reference was invalid, we terminate the process.



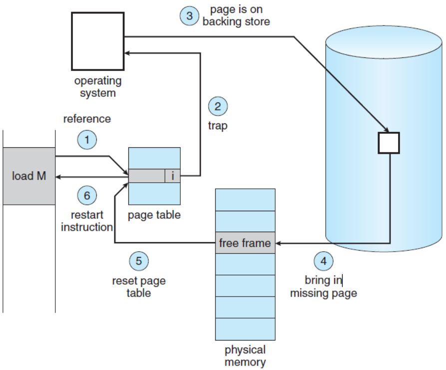
Handling Page Fault (2/6)

- If the reference was invalid, we terminate the process.
- If it was valid but we have not yet brought in that page, we now page it in.



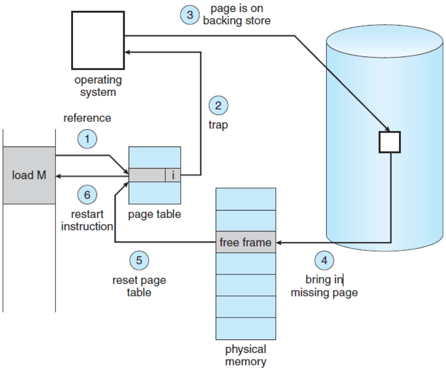
Handling Page Fault (3/6)

- We find a free frame.



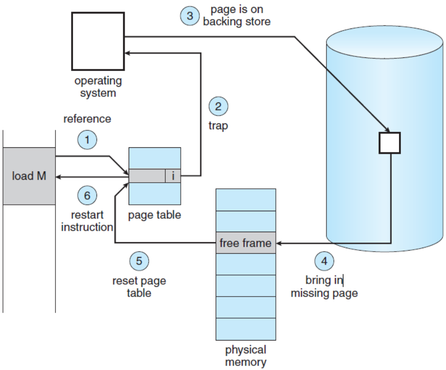
Handling Page Fault (4/6)

- We schedule a disk operation to read the desired page into the newly allocated frame.



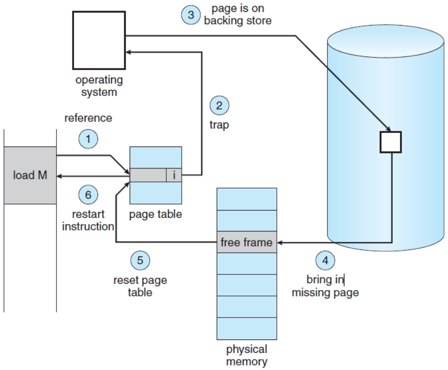
Handling Page Fault (5/6)

- When the disk read is complete, we modify the internal table kept with the process and the page table to indicate that the page is now in memory.



Handling Page Fault (6/6)

- We restart the instruction that was interrupted by the trap.



Page Replacement

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- And, we run 6 processes, each of which is 10 pages in size, but actually uses only 5 pages.

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- Increasing the degree of multiprogramming: over-allocating memory

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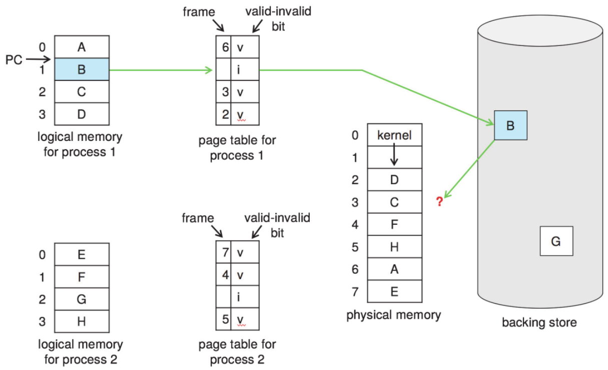
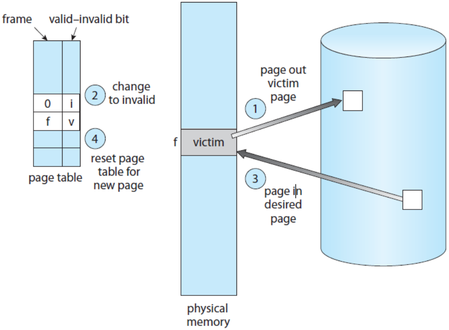
Need For Page Replacement



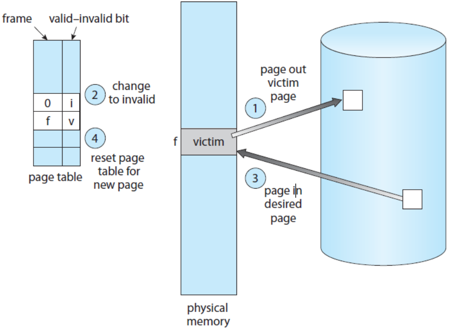
Page Replacement (1/4)

- Find the location of the desired page on disk.



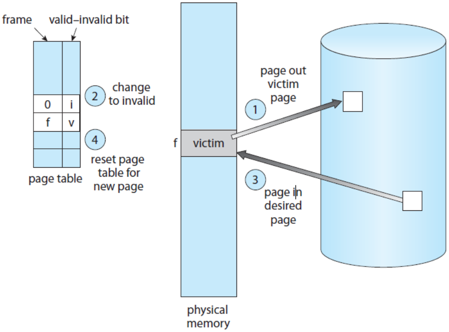
Page Replacement (2/4)

- Find a free frame.



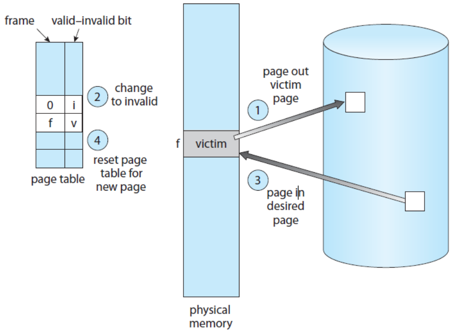
Page Replacement (2/4)

- Find a free frame.
  - If there is a free frame, use it.



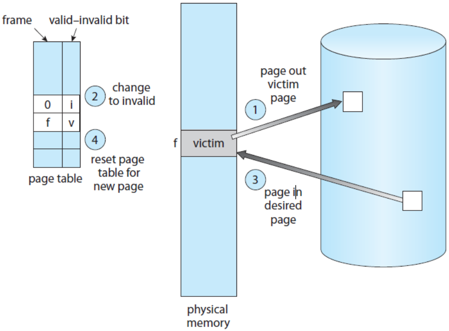
Page Replacement (2/4)

- Find a free frame.
  - If there is a free frame, use it.
  - If there is no free frame, use a page replacement algorithm to select a victim frame: write victim frame to disk if dirty.



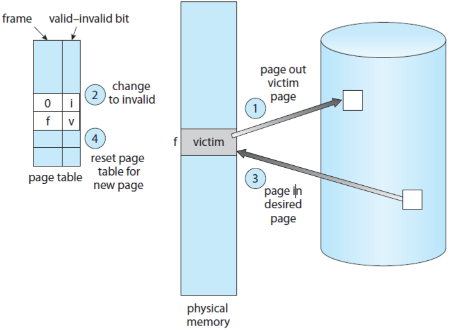
Page Replacement (3/4)

- Bring the desired page into the (newly) free frame; update the page and frame tables

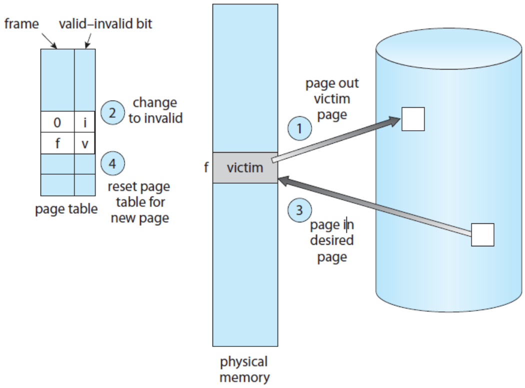


Page Replacement (4/4)

- Continue the process by restarting the instruction that caused the trap.



Dirty Bit



- Use modify (dirty) bit to reduce overhead of page transfers - only modified pages are written to disk.

Page Replacement Algorithms

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- Refernce string is a sequence of page numbers.

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Evaluate Page Replacement Algorithms

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  7, 0, 1, 2, 0, 3, 0, 4, 2, 3, 0, 3, 0, 3, 2, 1, 2, 0, 1, 7, 0, 1
- Evaluate algorithm by running it on a reference string and computing the number of page faults on that string.

Page Replacement Algorithms

- First-In-First-Out (FIFO)
- Optimal
- Least Recently Used (LRU)
- LRU-Approximation
- Counting-Based

FIFO Page Replacement

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- Reference string: 7,0,1,2,0,3,0,4,2,3,0,3,0,3,2,1,2,0,1,7,0,1

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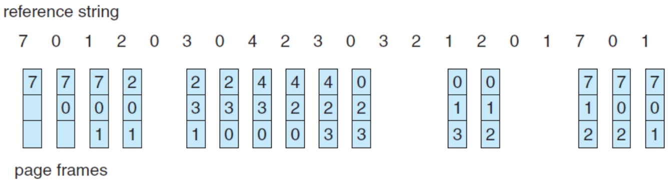

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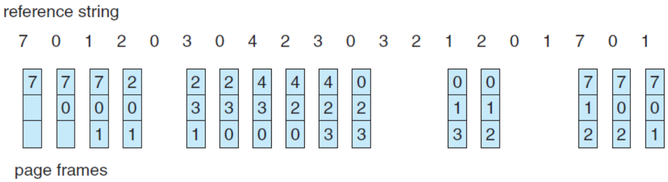

- 15 page faults

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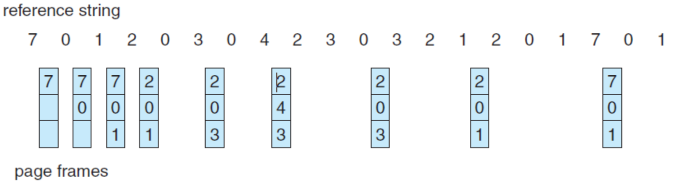

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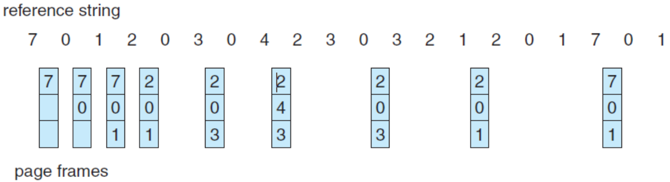

- Used for measuring how well your algorithm performs.

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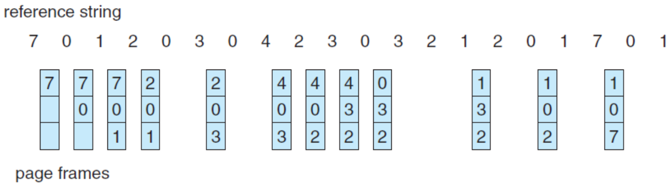

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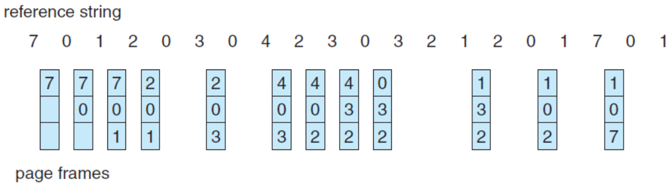

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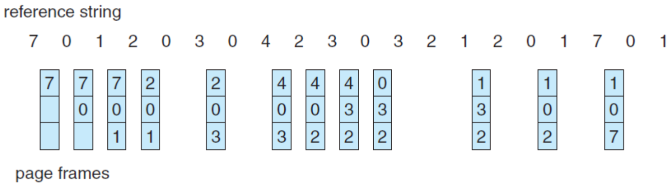

- 12 faults: better than FIFO but worse than OPT
- Generally good algorithm and frequently used

LRU Implementation (1/2)

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- Search through table needed.

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LRU Implementation (2/2)

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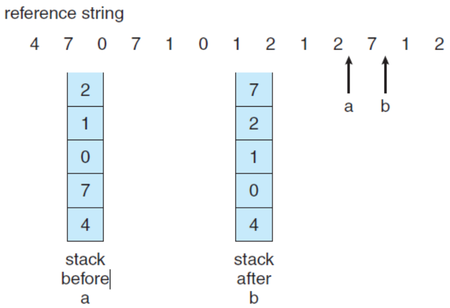
- Page referenced: move it to the top
- No search for replacement.

Stack Implementation

- Use of a stack to record most recent page references.



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  - Second-chance algorithm

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- We do not know the order

Second-Chance Algorithm (1/2)

- Generally FIFO, plus hardware-provided reference bit

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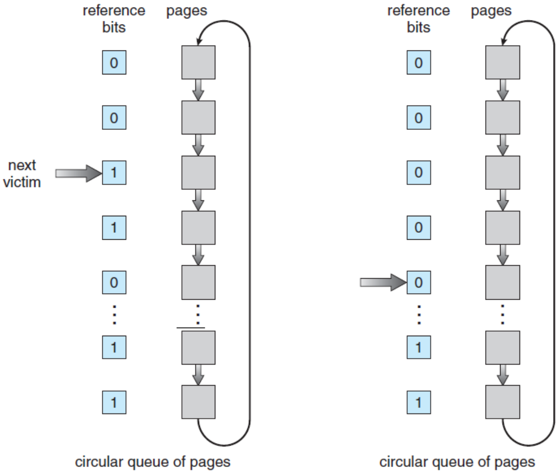
- If page to be replaced has
  - Reference bit $= 0 \rightarrow$ replace it
  - Reference bit $= 1$ then, set reference bit 0, leave page in memory, and replace next page, subject to same rules.

Second-Chance Algorithm (2/2)



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- Lease Frequently Used (LFU) algorithm: replaces page with smallest count.
- Most Frequently Used (MFU) algorithm: based on the argument that the page with the smallest count was probably just brought in and has yet to be used.

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- Demand paging similar to paging + swapping
- Page fault
- Page replacement algorithms:
  - FIFO, optimal, LRU, LRU-approximate, counting-based

Questions?

Acknowledgements

Some slides were derived from Avi Silberschatz slides.